Subgame perfect equilibria in majoritarian bargaining

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ABSTRACT

We study the division of a surplus under majoritarian bargaining in the three-person case. In a stationary equilibrium as derived by Baron and Ferejohn (1989), the proposer offers one third times the discount factor of the surplus to a second player and allocates no payoff to the third player, a proposal which is accepted without delay. Laboratory experiments show various deviations from this equilibrium, where different offers are typically made and delay may occur before acceptance. We address the issue to what extent these findings are compatible with subgame perfect equilibrium and characterize the set of subgame perfect equilibrium payoffs for any value of the discount factor. We show that for any proposal in the interior of the space of possible agreements there exists a discount factor such that the proposal is made and accepted. We characterize the values of the discount factor for which equilibrium with one-period delay exist. We show that any amount of equilibrium delay is possible and we construct subgame perfect equilibria such that arbitrary long delay occurs with probability one.

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1. Introduction

Many decisions in legislatures involve the task of allocating resources among constituencies with diverse and often opposing preferences. To analyze such decisions, many authors have studied extensions of the Rubinstein (1982) bargaining model to collective choice problems, see for instance Baron and Ferejohn (1989), Harrington (1990), Baron and Kalai (1993), and Banks and Duggan (2000). These papers study multilateral bargaining games, where an alternative is accepted if it is approved by a set of players that belongs to a collection of decisive coalitions. This approach makes it possible to study a variety of important institutional set-ups, including the one of majority voting.

The Baron and Ferejohn (1989) model, hereafter BF model, is the most frequently used model to study legislative bargaining. Baron and Ferejohn (1989) consider a bargaining game with infinite time horizon, where the proposer selection process is modeled by time-invariant recognition probabilities, and acceptance by a simple majority is sufficient to implement a proposal. Eraslan (2002) shows that stationary subgame perfect equilibrium payoffs are unique for a rather general specification of the BF model, and Eraslan and McLennan (2013) show uniqueness of such equilibrium payoffs in an even more general model. In any stationary equilibrium of the BF model, the proposer offers $\delta/n$ to half of the responding players, nothing to the other responders, and keeps the remainder. Such a proposal is accepted without delay. The equilibrium therefore predicts substantial proposer power, the formation of a minimum winning coalition, and absence of delay in reaching an agreement.

The BF result has been used as the theoretical benchmark in many papers in the experimental literature. McKelvey (1991) studies a bargaining game with three voters and a discrete policy space consisting of three alternatives. Fréchette et al. (2003) use the BF model with five players as their main treatment and study the effects of open versus closed amendment rules. Fréchette et al. (2005) use the BF specification with three players and a discount factor of $1/2$ in their treatments, as well as the case without discounting, and focus on the influence of nominal bargaining power on bargaining outcomes. Diermeier and Morton (2006) consider a version of the BF model with a finite time horizon and heterogeneous recognition probabilities. Breitmoser and Tan (2010) consider the BF specification with three players in their treatments and study in particular whether fairness can explain some of the experimental findings. Kagel et al. (2010) are interested in the consequences of veto power in committees and use the BF specification with three players and a discount factor of 0.95 as a benchmark. Miller and Vanberg (2013) use the BF model with three players and a discount factor of 0.9 when comparing the costs of reaching agreement under majority and unanimity rule. Agranov and Tergiman (2014) and Baranski and Kagel (2015) allow players to communicate after the proposer is selected but before a proposal is submitted and respectively study the BF model with five players and a discount factor of 0.8 and the BF model with three players and no discounting.

The experimental literature finds that the predictions of the stationary subgame perfect equilibrium in the BF model are violated.
in several ways, although certain qualitative features like the presence of proposer power and a tendency to form minimum winning coalitions are found. The most striking feature of the data is that proposers make overly generous offers to the responding players and keep a much smaller share for themselves than predicted by the stationary equilibrium. Minimum winning coalitions are not always formed and even an equal sharing of the surplus with all players occurs with significantly positive probability. Contrary to the BF prediction of immediate acceptance, there is a significant number of observations with delay before a proposal is accepted, in particular when the discount factor is sufficiently high.

McKelvey (1991) conjectures that subjects may offer too much to potential coalition partners because of fear of retaliation in later rounds and Diermeier and Morton (2006) report the use of non-stationary strategies. For instance, subjects who vote to reject a proposal on average receive a higher payoff from the new proposal, which is a violation of stationarity since under stationarity irrelevant parts of the history should be of no influence to future behavior. Bradfield and Kagel (2015) experimentally investigate the BF model when the actual bargaining is done by teams, as is usually the case in committee bargaining, and find clear evidence for teams retaliating more against proposers whose proposals were rejected. Many authors, most notably Rubinstein (1991), have argued against the use of stationarity strategies in bargaining. We are therefore interested in the question whether proposals like for instance equal sharing of the surplus are compatible with subgame perfect equilibrium and whether delay is possible in subgame perfect equilibria when stationarity is not imposed.

In this paper we study the BF model with three players. Since many parliamentary democracies involve three major parties, the three player case is of considerable importance. We characterize the set of subgame perfect equilibrium utilities for any value of the discount factor. When the discount factor goes to zero, the set of subgame equilibrium payoffs shrinks to the unit vector and the entire surplus is allocated to the player who is recognized as the proposer in period zero. When the discount factor goes to one, the set of subgame equilibria payoffs expands to the set of feasible payoffs. In particular, all inefficient outcomes are supported as subgame perfect equilibria when the discount factor is sufficiently high. Our characterization of the set of subgame perfect equilibrium utilities is entirely constructive.

Subgame perfect equilibrium, as a function of the discount factor, exhibits different qualitative properties in the three player BF model than in the BF model with five or more players. Baron and Ferejohn (1989) show that in the BF model with five players any feasible payoff vector may be supported as a subgame perfect equilibrium payoff for a discount factor exceeding 7/8. Thus in particular, even the extreme allocations corresponding to the vertices of the set of feasible payoffs could be proposed and accepted in a subgame perfect equilibrium. Likewise, in the BF model with more than five players, there exists a discount factor such that for any feasible division of the surplus there is a subgame perfect equilibrium such that the given division is proposed in round 0 and accepted.

This result is not true for the three player BF model. For any given value of the discount factor, every player’s equilibrium payoff has a positive lower bound. It follows that the allocation where one player gets the entire cake can never be supported as a subgame perfect equilibrium outcome, irrespectively of the discount factor. Furthermore, for any given value of the discount factor, payoff vectors close to the boundary of the feasible set cannot be supported as subgame perfect equilibrium outcomes. It is nevertheless the case that for any vector in the interior of the set of feasible payoffs there exists a discount factor such that this vector is proposed and accepted in period 0 in a subgame perfect equilibrium.

The results on unanimity bargaining games with three or more players are likewise different from ours. Thus for instance in Shaked’s example of unanimity bargaining with three players as presented in Binmore et al. (1992), for a discount factor greater than 1/2, any given feasible payoff vector could be proposed and accepted in a subgame perfect equilibrium. In contrast, in the three player BF model studied here, for any discount factor greater than 1/2, player’s equilibrium payoff is bounded below by 0.26.

Our paper complements the findings in Li (2014), who also studies the three player BF model. First of all, we consider the entire set of subgame perfect equilibrium payoffs, both efficient and inefficient, while Li (2014) focuses on the set of efficient proposals. Inefficiency can arise both because inefficient divisions of the surplus could be proposed and accepted and because of delay. Our analysis captures both types of inefficiencies. Thus in particular we derive the range of discount factors under which the model admits an equilibrium with one-period delay. Secondly, Li (2014) considers a version of the model with simultaneous voting, whereas we take the version where voting is sequential. As a consequence, our results do not rely on any additional refinement of subgame perfection such as the requirement that the players do not use weakly dominated strategies. Finally, the result of Li (2014) relies on mixed strategies, whereas we only use pure strategies. It turns out that the set of efficient payoff vectors that can be supported as a subgame perfect equilibrium payoff in mixed strategies can also be sustained using pure strategies.

Chatterjee et al. (1993) consider bargaining in the context of strictly superadditive transferable utility games under the rejector-proposes protocol and show the usual statement of the folk theorem that for a sufficiently high discount factor any individually rational outcome is possible. Note that the BF model violates strict superadditivity when viewed as a transferable utility game. Norman (2002) studies a finite horizon version of the BF model and shows for the case with at least five players that any interior division of the surplus can be supported as a subgame perfect equilibrium outcome if players are sufficiently patient and there are sufficiently many rounds of bargaining. He also shows indeterminacy of the equilibrium outcome whenever there are at least three players and at least three bargaining rounds.

The BF model with three players is a borderline case where the usual statement of the folk theorem does not apply, but the set of subgame perfect equilibria can be sizable. Recall that in the Rubinstein (1982) model with two players, there is a unique subgame perfect equilibrium. Cho and Duggan (2009) analyze a version of the BF model where an arbitrary number of players bargains on a one-dimensional ideological spectrum of policies and agents have single-peaked preferences over policies. Under sequential voting and some mild assumptions on the order of voting, they find that only the median ideal alternative can be supported in a subgame perfect equilibrium. Cho and Duggan (2015) demonstrate that this result does not carry over to simultaneous and publicly observed voting, even when voting is assumed to be stage-undominated, and obtain the support of arbitrary outcomes for arbitrary positive discount factors.

For the BF model with three players, we provide necessary and sufficient conditions for one period of delay before acceptance takes place. We characterize the threshold such that, for any value of the discount factor greater than the threshold, one period of delay is possible and for any value of the discount factor smaller than the threshold all subgame perfect equilibria are characterized by immediate acceptance. We also demonstrate the possibility of arbitrary long delay before acceptance takes place in a subgame perfect equilibrium when the discount factor is sufficiently high.

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1 For example, the British political system of the twentieth century consisted of three major parties representing the Conservatives, Labor, and Liberals. Similarly, for a long time the German political system had three major parties, the Christian Democrats, the Social Democrats, and the Liberal Party.
For low values of the discount factor, the offers as observed in experiments are not consistent with stationary equilibrium, and are not even consistent with subgame perfect equilibrium. For instance, when the discount factor is equal to 1/2, Fréchette et al. (2005) find an average share of the proposer for accepted offers equal to 0.50, lower than the lowest possible share of 0.78 in a subgame perfect equilibrium, which does not differ much from the stationary equilibrium share of 0.83. For higher values of the discount factor, the literature reports on experiments with discount factors equal to 0.9, 0.95, and 1, the experimental findings on accepted offers are not in contradiction with the predictions of subgame perfect equilibrium. The experiments by Agranov and Tergiman (2014) and Baranski and Kagel (2015) allow for communication between the players after the proposer is determined but before a proposal is made. In treatments with communication, proposals are much closer to those of the stationary equilibrium. This suggests that communication can play a role in the selection of a particular equilibrium.

Regarding delay before reaching an agreement, the experimental literature finds little delay for experienced subjects when the discount factor is equal to 0.5, but substantial delay when the discount factor is equal to 0.9 or higher. We find that the critical discount factor needed to have subgame perfect equilibria with one-period delay is slightly above 0.8.

The rest of the paper is organized as follows. In Section 2 we describe the BF model and derive bounds on subgame perfect equilibrium utilities as a function of the discount factor. In Section 3 we characterize the subgame perfect equilibrium utilities conditional on the selection of a proposer and argue that the bounds in Section 2 are tight. The explicit construction of the subgame perfect equilibria that reach the bounds is relegated to the Appendix. In Section 4 we provide necessary and sufficient conditions on the discount factor for one period of delay before acceptance takes place. In Section 5 we investigate the possibility of arbitrary long delay before acceptance takes place in a subgame perfect equilibrium. Section 6 concludes.

2. The bounds for subgame perfect equilibrium payoffs

We consider the three player version of the BF model. The set of players \( N = \{1, 2, 3\} \) has to agree on the choice of a payoff vector in the set of feasible payoffs

\[
V = \{x \in \mathbb{R}_+^3 \mid x_1 + x_2 + x_3 \leq 1\}.
\]

In each time period \( t = 0, 1, 2, \ldots \) nature selects a proposer and the order of the responders by means of time-invariant recognition probabilities. More precisely, in the beginning of period \( t = 0, 1, 2, \ldots \) nature selects each permutation \( \pi^t : \{1, 2, 3\} \to N \) in the set of six possible permutations \( I^t \) with equal probability. The choice of the permutation is independent across time periods. Player \( \pi^t(1) \) makes a proposal \( x^t \in V \). Next Player \( \pi^t(2) \) responds by accepting or rejecting the proposal. If \( \pi^t(2) \) accepts the proposal, the game ends, and the proposal \( x^t \) is implemented. Otherwise, Player \( \pi^t(3) \) reacts to the proposal. If Player \( \pi^t(3) \) accepts the proposal, the game ends. If Player \( \pi^t(3) \) rejects the proposal, period \( t + 1 \) begins. The utility of player \( i \in N \) who receives outcome \( x_i \) in period \( t \) is \( \delta^t x_i \), where \( \delta \in [0, 1) \) is the common discount factor. If no agreement is ever reached, payoffs are 0.

A history \( h \) is the sequence of all actions that have occurred before a particular decision node in the game. For simplicity we suppress the elements of the sequence pertaining to the moves by the responding players: it is understood that in any non-terminal history both responders have rejected all proposals to date. With this convention, for \( t = 0, 1, \ldots \) any non-terminal history is of one of the following two types:

1. \( h \in H^2 \), if and only if \( h \) is of the form \( (\pi^0, x^0, \ldots, x^{t-1}, \pi^t) \).
2. \( h \in H^3 \), if and only if \( h \) is of the form \( (\pi^0, x^0, \ldots, x^{t-1}, \pi^t, x^t) \),

where for every \( k \in \{0, \ldots, t\} \), \( \pi^k \) is a permutation of \( N \) and \( x^k \in V \). After a history \( h \in H^2 \), the proposer makes a move and, after a history \( h \in H^3 \), a responder accepts or rejects the current proposal. Histories in \( H^2 \) are therefore called proposer histories and those in \( H^3 \) responder histories.

A pure strategy for Player \( i \in N \) assigns to every history \( h = (\pi^0, x^0, \ldots, \pi^t) \in H^2 \) with \( \pi^t(1) = i \) an element of \( V \), and to every history \( h \in H^3 \) with \( \pi^t(1) \neq i \) an accept/reject decision. All our results make only use of pure strategy profiles. It is an easy exercise to verify that the bounds we present on equilibrium utilities carry over to mixed strategy profiles as well.

A strategy is stationary if it does not depend on payoff irrelevant parts of the history, see Maskin and Tirole (2001) for a precise formulation of the notion of stationarity. For the BF model, stationarity means that a player makes the same proposal at every history in \( H^2 \) where he is recognized as a proposer and bases his accept/reject decisions only on the current proposal and the identity of the player voting next to him in the current voting round.

The following result follows from the analysis in Baron and Ferejohn (1989).

**Theorem 2.1.** For every \( \delta \in [0, 1) \), a strategy profile is a subgame perfect equilibrium in stationary strategies if and only if it has the following form: (1) the recognized player offers \( \delta/3 \) to one responding player and keeps \( 1 - \delta/3 \) for himself; (2) a responding player accepts any proposal in which he receives at least \( \delta/3 \) and rejects otherwise; (3) all players receive \( \delta/3 \) as responding player with ex ante probability equal to \( 1/3 \).

In equilibrium, the recognized player forms a minimum winning coalition with one responding player, offers the responding player exactly his continuation utility, offers nothing to the third player, and keeps the remainder of the surplus. The responding player who is part of the minimum winning coalition accepts this proposal.\(^2\) The equilibrium in Theorem 2.1 is not unique, since there is some indeterminacy in the selection by the proposer of the player who joins the minimum winning coalition. However, all strategy profiles described in Theorem 2.1 lead to the same equilibrium payoffs. Baron and Ferejohn (1989) show that this result extends to the case with an odd number \( n \) of players, where now \( \delta(n - 1)/2n \) is offered to \( (n - 1)/2 \) players, 0 to the other players, and the remainder is kept by the proposer. It follows from Eraslan (2002) that stationary subgame perfect equilibrium payoffs are unique, also when mixed strategies are allowed for.

We now drop the stationarity assumption on strategies. The set of all subgame perfect equilibrium (SPE) payoff vectors is denoted by \( U \), a subset of \( V \). For \( i \in N \), let \( u_i^+ \) be the supremum of SPE payoffs for Player \( i \) and \( u_i^- \) be the infimum, so we want to show that

\[
u_i^+ = \sup\{u_i \in \mathbb{R} \mid u \in U\},
\]

\[
u_i^- = \inf\{u_i \in \mathbb{R} \mid u \in U\}.
\]

We define the real numbers \( \underline{b} \) and \( \overline{b} \) by

\[
\underline{b} = \frac{3 - 3\delta}{\delta - 6\delta - \delta^2},
\]

\[
\overline{b} = \frac{3 - \delta}{\delta - 6\delta - \delta^2}.
\]
The statement in [1] of Claim 2.3 is subtle. It considers a proposal where one or both responding players receive more than $\delta \overline{u}$. The claim is that such a proposal is accepted, but the statement does not specify by whom. Indeed, if both responding players receive more than $\delta \overline{u}$, then Player $\pi(2)$ can safely reject such a proposal in the knowledge that Player $\pi(3)$ is going to accept it. Similarly, if the proposal is such that only Player $\pi(3)$ receives more than $\delta \overline{u}$, it could still be that Player $\pi(2)$ accepts it, since Player $\pi(2)$ is indifferent whether such a proposal is accepted by himself or by Player $\pi(3)$.

The following claim derives bounds on SPE utilities conditional on the selection of the proposer.

Claim 2.4. Consider the subgame following the period 0 proposer history $h = (\pi)$ and let $v = (v_1, v_2, v_3)$ be the payoffs of an SPE $s$ conditional on the subgame being reached. It holds that

1. $v_{\pi(1)}(1) \geq 1 - \delta \overline{u}$.
2. $v_{\pi(1)}(1) \leq 1 - \delta \overline{u}$.
3. $v_{\pi(2)}(2) \leq \delta \overline{u}$ and $v_{\pi(3)}(3) \leq \delta \overline{u}$.

Proof. [1] By Claim 2.3 the proposal $x$ where $x_{\pi(1)}(1) = 1 - \delta \overline{u} - \epsilon$, $x_{\pi(2)}(2) = \delta \overline{u} + \epsilon$, and $x_{\pi(3)}(3) = 0$ is accepted for each $\epsilon > 0$. Since $s$ is an SPE we have $v_{\pi(1)}(1) \geq 1 - \delta \overline{u} - \epsilon$ for each $\epsilon > 0$. It follows that $v_{\pi(1)}(1) \geq 1 - \delta \overline{u}$.

[2] Let $x$ be Player $\pi(1)$’s proposal at $h$ under the strategy profile $s$. Suppose first that $x$ is rejected under $s$. Since this result holds at least one period of delay, it follows that $v_{\pi(1)}(1) \leq \delta \overline{u}$. It follows from Claim 2.2 that $\delta \overline{u} \leq 1 - \delta u$ and therefore that $v_{\pi(1)}(1) \leq 1 - \delta u$. Suppose now $x$ is accepted. Then by Claim 2.3 it is the case that $x_{\pi(2)}(2) \geq \delta u$ or $x_{\pi(3)}(3) \geq \delta u$. Therefore $v_{\pi(1)}(1) = x_{\pi(1)}(1) \leq 1 - x_{\pi(2)}(2) - x_{\pi(3)}(3) \leq 1 - \delta u$.

[3] We have $v_{\pi(2)}(2) \leq 1 - v_{\pi(1)}(1) - v_{\pi(3)}(3) \leq 1 - v_{\pi(1)}(1) \leq \delta \overline{u}$, where the last inequality follows from [1]. The argument for $v_{\pi(3)}(3)$ is similar. \hfill \Box

The equilibrium utility conditional on being the initial proposer is in between $1 - \delta \overline{u}$ and $1 - \delta u$. Responders have equilibrium utilities bounded from above by $\delta \overline{u}$. The result carries over to subgames in later periods $t$, except that discounting implies that all utilities should be multiplied by $\delta^t$.

We now use the claims derived so far to derive a lower bound on $u$ in terms of $\delta$ and $\overline{u}$, and an upper bound on $\overline{u}$ in terms of $\delta$ and $u$.

Claim 2.5. It holds that

\[
\frac{1 - \delta \overline{u}}{3} \leq u \leq \overline{u} \leq \frac{1 - \delta u}{3 - 2\delta}.
\]

Proof. Let $s$ be an SPE with payoffs $u \in U$. For each permutation $\pi \in \Pi$, let $v' = (v'_1, v'_2, v'_3)$ be the payoff of $s$ conditional on nature choosing $\pi$ as the initial permutation. Since each permutation is equally likely, we have

\[
u = \frac{1}{3} \sum_{\pi \in \Pi} v'.
\]

For a permutation $\pi \in \{(1, 2, 3), (1, 3, 2)\}$ it holds that Player 1 is the proposer. By parts [1] and [2] of Claim 2.4 it holds that

\[
1 - \delta \overline{u} \leq v'_1 \leq 1 - \delta u.
\]
Next, consider a permutation \( \pi \in \Pi \setminus \{(1, 2, 3), (1, 3, 2)\} \). Now Player 1 is a responding player and so \( 0 \leq v_1^* \leq \delta \bar{u} \) by part [3] of Claim 2.4. We conclude that
\[
\frac{2}{6}(1 - \delta \bar{u}) + \frac{4}{6}0 \leq u_1 \leq \frac{2}{6}(1 - \delta u) + \frac{4}{6} \delta \bar{u}.
\]
Since the inequalities hold for every \( u \in U \), we have
\[
\frac{1}{3}(1 - \delta \bar{u}) \leq u \leq \frac{1}{3}(1 - \delta u) + \frac{2}{3} \delta \bar{u}.
\]
Rearranging the last inequality yields the result. \( \square \)

We are now in a position to prove that the payoff to any player in any SPE of the game is bounded from below by \( \bar{b} \) and from above by \( \bar{b} \).

**Theorem 2.6.** It holds that \( \bar{b} \leq u \leq \bar{u} \leq \bar{b} \).

**Proof.** Define the functions \( f, g : [0, 1] \rightarrow [0, 1] \) by
\[
f(y) = \frac{1 - \delta y}{3} \text{ and } g(y) = \frac{1 - \delta y}{3 - 2 \delta},
\]
and let \( h = g \circ f \). Then the function \( h \) is defined by
\[
h(y) = \frac{1 - \delta (\frac{1}{3} - \delta y)}{3 - 2 \delta} = \frac{3 - \delta + \delta^2 y}{9 - 6 \delta}, \quad y \in [0, 1].
\]
By Claim 2.5, we have that
\[
f(\bar{u}) \leq u \leq \bar{u} \leq g(u).
\]
Using the fact that \( g \) is a non-increasing function, we obtain \( \bar{u} \leq g(f(\bar{u})) \), and hence \( \bar{u} \leq h(\bar{u}) \). Since \( h \) is a non-decreasing function, we can iterate the last inequality to obtain \( \bar{u} \leq h^n(\bar{u}) \).

Both \( f \) and \( g \) are contractions, since \( \frac{1}{3} < 1 \) and \( \bar{u} > \bar{b} \). Hence \( h \) is also a contraction. Moreover, \( \bar{b} \) is a fixed point of \( h \). By the Banach contraction theorem it follows that \( h^n(\bar{u}) \) converges to \( \bar{b} \), so we have shown that \( \bar{u} \leq \bar{b} \).

An easy computation shows that \( \bar{b} = f(\bar{b}) \). Since \( f \) is a non-increasing function, we have that
\[
\bar{b} = f(\bar{b}) \leq f(\bar{u}) \leq u,
\]
which completes the proof. \( \square \)

### 3. Explicit subgame perfect equilibria

In the previous section we have shown that SPE utilities are bounded from below by \( \bar{b} \) and from above by \( \bar{b} \). In this section we show that these bounds are tight by constructing an SPE yielding Player 2 a payoff of \( \bar{b} \) and Player 3 a payoff of \( b \).

We define
\[
V_1 = \{ v \in V \mid v_1 \geq 1 - \delta \bar{b} \text{ and } v_2 \geq \delta \bar{b} \text{ or } v_3 \geq \delta \bar{b} \},
\]
\[
V_2 = \{ v \in V \mid v_2 \geq 1 - \delta \bar{b} \text{ and } v_3 \geq \delta \bar{b} \text{ or } v_1 \geq \delta \bar{b} \},
\]
\[
V_3 = \{ v \in V \mid v_3 \geq 1 - \delta \bar{b} \text{ and } v_1 \geq \delta \bar{b} \text{ or } v_2 \geq \delta \bar{b} \}.
\]

For \( i \in N \), the set \( V_i \) gives the payoffs that can be obtained as SPE utilities conditional on reaching a subgame in period 0 where Player \( i \) is selected as a proposer.

**Theorem 3.1.** Choose \( a_1 \in V_1, a_2 \in V_2, \) and \( a_3 \in V_3 \). Then there exists an SPE such that for each permutation \( \pi^0 \) the proposal \( a_{\pi^0(1)} \) is made and accepted in period 0.

The constructive proof of Theorem 3.1 can be found in the Appendix. Using Theorem 3.1, it is easy to prove our main result.

**Theorem 3.2.** It holds that \( \bar{u}_1 = \bar{u}_2 = \bar{u}_3 = \bar{b} \) and \( u_1 = u_2 = u_3 = b \).

**Proof.** For \( i \in N \), we choose elements \( a_i \in V_i \) that are the best for Player 2 and the worst for Player 3, so \( a_1 = (1 - \delta \bar{b}, \delta \bar{b}, 0), \) \( a_2 = (\delta \bar{b}, 1 - \delta \bar{b}, 0) \), and \( a_3 = (0, \delta \bar{b}, 1 - \delta \bar{b}) \). Now let \( s \) be a strategy profile that satisfies the conditions of Theorem 3.1. Since
\[
\bar{b} = \frac{1}{3}(1 - \delta \bar{b}),
\]
\[
\bar{b} = \frac{1}{3}(1 - \delta \bar{b}) + \frac{2}{3} \delta \bar{b},
\]
we find that the strategy profile \( s \) yields a payoff of \( \bar{b} \) for Player 2 and a payoff of \( \bar{b} \) for Player 3. \( \square \)

The example below illustrates that an inefficient proposal could be proposed and accepted in period 0.

**Example 3.3.** Take \( \delta = 1/3 \), so \( \bar{b} = 9/31 \) and \( \bar{b} = 12/31 \). We have that \( 1 - \delta \bar{b} = 27/31 \) and \( \delta \bar{b} = 3/31 \). Let \( a_1 = (27/31, 3/31, 0), a_2 = (0, 27/31, 3/31), \) and \( a_3 = (0, 3/31, 0, 27/31) \). As \( a_1, a_2, \) and \( a_3 \) lie in the sets \( V_1, V_2, \) and \( V_3 \), respectively, there is a subgame perfect equilibrium such that, depending on the identity of the proposer, the points \( a_1, a_2, \) and \( a_3 \) are proposed and accepted. Each of the proposals is inefficient.

It is illuminating to compare our findings here with folk theorems for unanimity bargaining. Consider a game with primitives as in Section 2 where the proposal is implemented only if approved of by both responders. The main result in Herings et al. (2017) applies to show that for each \( \delta > 1/2 \) and each point \( u \) in \( V \) there exists a subgame perfect equilibrium such that \( u \) is proposed (no matter who the proposer is) and is approved of by both responders in period 0. Thus all the payoff vectors in \( V \) could be supported as subgame perfect equilibrium payoffs. Related folk theorems for unanimity bargaining games have been established in e.g. Herrero (1985) and Haller (1986). In contrast, in the majoritarian bargaining game, for any fixed discount factor, there exist payoff vectors in \( V \) that cannot be supported as a subgame perfect equilibrium. This shows that the result for the unanimity case is qualitatively different from the majoritarian case.

In any SPE in stationary strategies, the proposer obtains a payoff of \( 1 - \delta/3 \) in period 0 and one of the responders receives a payoff of \( \delta/3 \). The stationary equilibrium is usually taken as a theoretical benchmark in the experimental literature. In this literature, substantial deviations from this benchmark are found. McKelvey (1991) describes an experiment with a discrete version of the BF model. Compared to the stationary prediction, proposers usually offer too much to the responding players. Diermeier and Morton (2006) consider a finitely repeated version of the BF model. Similar to McKelvey (1991) and Diermeier and Morton (2006) find little support for the predictions of the stationary equilibrium. First, in one third of the cases a positive amount was proposed to all players, not just to the members of the minimal winning coalition. Moreover, proposers consistently offer too much to other coalition members. Both McKelvey (1991) and Diermeier and Morton (2006) report the use of non-stationary strategies. We verify to what extent the findings in the experimental literature are consistent with the non-stationary subgame perfect equilibria as found in Theorem 3.1. We refer to the responder who is offered the higher share in the cake as Player 2 and to the other responder as Player 1.

Fréchette et al. (2005) provide experimental evidence for the under-realization of proposer power and the generosity shown to voters outside the minimal winning coalition for the BF model with \( \delta = 0.5 \) and \( \delta = 1 \). They consider both treatments with experienced and treatments with inexperienced players.
When $\delta = 1$, 39% of the proposals for inexperienced players do not correspond to the formation of a minimum winning coalition, a number that is still equal to 23% for experienced players. Even perfectly egalitarian proposals are observed in more than 5% of the cases. A perfectly egalitarian proposal is consistent with subgame perfect equilibrium when $1 - \delta \delta_b \leq 1/3$, which is equivalent to $\delta \geq (21 - 3/\sqrt{2})/2 \approx 0.895$. The average share of the proposer for accepted offers is 0.51 for inexperienced players and 0.52 for experienced players, considerably lower than the stationary equilibrium prediction of 0.67. The average share of Player $j$ is 0.43 for inexperienced players and 0.45 for experienced players, considerably more than the stationary equilibrium prediction of 0.33. Since we have shown that any proposal and any pay off is possible in an SPE for discount factors tending to one, these results would not contradict SPE.

When $\delta = 0.5$, 57% of the proposals for inexperienced players do not correspond to the formation of a minimum winning coalition, a number that is still equal to 23% for experienced players. The average share of the proposer for accepted offers is 0.50 for inexperienced players and 0.59 for experienced players, considerably lower than the stationary equilibrium prediction of 0.83. The average share of Player $j$ is 0.41 for inexperienced players and 0.39 for experienced players, considerably more than the stationary equilibrium prediction of 0.17. Our results predict a payoff to the proposer in the range from $1 - \delta \delta_b = 0.78$ to $1 - \delta \delta_b = 0.87$. The payoff to the player who receives the higher share is predicted to be in the range from $\delta \delta_b = 0.13$ to $\delta \delta_b = 0.22$. The payoff to the player who receives the lower share is predicted to be in the range from 0 to $\delta \delta_b/2 = 0.11$. For $\delta = 0.5$ the experimental findings are not even consistent with SPE. This is not so surprising, since for such low values of the discount factor, the potential influence of punishment is quite limited. Fréchette et al. (2005) offer the reluctance of coalition partners to accept offers much below 1/3 as an explanation of these findings.

Kagel et al. (2010) consider a treatment corresponding to the BF model with a discount factor $\delta = 0.95$. They report that 40.8% of the proposals do not correspond to the formation of a minimum winning coalition. For accepted offers, the average share of the proposer is 0.458 and for Player $j$ it is 0.44. Very similar results are found in Breitmoser and Tan (2010), who also study the case $\delta = 0.95$. The sample estimates are 0.445 to the proposer, 0.359 to Player $j$, and 0.172 to Player $k$. The stationary equilibrium predicts a payoff of 0.683 to the proposer, 0.317 to Player $j$, and 0 to Player $k$. Our results predict a payoff to the proposer in the range from $1 - \delta \delta_b = 0.188$ to $1 - \delta \delta_b = 0.938$. The payoff to the player who receives the higher share is predicted to be in the range from $\delta \delta_b = 0.063$ to $\delta \delta_b = 0.813$. The payoff to the player who receives the lower share is predicted to be in the range from 0 to $\delta \delta_b = 0.406$.

Miller and Vanberg (2013) compare the costs of reaching agreement under majority and unanimity rule in the context of an experimental bargaining game with $\delta = 0.9$. They find patterns very similar to those reported in the previous literature. Proposers demand a higher share than they allocate to non-proposers, but the difference is far from the equilibrium prediction. Interestingly, approximately half of the proposals in the first period are three-way equal splits and only one out of five allocates 0 to one of the non-proposers.

4. Equilibria with one-period delay

The experimental literature shows that one cannot rule out delay before acceptance takes place. Fréchette et al. (2005) conclude that a majority of the proposals are accepted without delay, however, delays persist until the end of the sessions. They find one or more rounds of delay for experienced subjects in 5% of the cases when $\delta = 0.5$, a number increasing to 23% when $\delta = 1$. For inexperienced subjects, one or more rounds of delay occur in 11% of the cases when $\delta = 0.5$ and in 32% of the cases when $\delta = 1$. Kagel et al. (2010) report one or more rounds of delay in 28% of the cases when $\delta = 0.95$. Miller and Vanberg (2013) consider the case where $\delta = 0.9$ and find one or more rounds of delay before acceptance takes place in 25% of the cases.

In this section we derive the necessary and sufficient conditions to have exactly one period of delay before acceptance takes place in an SPE. A strategy profile $s$ is said to have one-period delay if, irrespectively of the moves of nature, the proposal in period 0 is rejected by both responders, and the proposal in period 1 is accepted by at least one responder. More precisely, the strategy profile $s$ has one-period delay if for all permutations $\pi^0$ and $\pi^1$ the proposal $s_{\pi^0(1)}(\pi^0)$ is rejected by both Players $\pi^0(2)$ and $\pi^0(3)$, while the proposal $s_{\pi^1(1)}(\pi^0)$, $\pi^1(1)$ is accepted by $\pi^1(2)$ or $\pi^1(3)$.

Theorem 4.1. There exists an SPE with one-period delay if and only if $\delta \geq 6 - 3\sqrt{3} \approx 0.804$.

The condition $\delta \geq 6 - 3\sqrt{3}$ is derived from the inequality $1 - \delta \delta_b \leq \delta \delta_b$, which expresses that the minimum equilibrium payoff conditional on being the proposer should be less than or equal to the discounted maximum unconditional equilibrium payoff. We first prove the only if part of Theorem 4.1.

Claim 4.2. If there is an SPE with one-period delay, then $\delta \geq 6 - 3\sqrt{3}$.

Proof. Let $s$ be an SPE with one-period delay. Let $\pi^0 = (1, 2, 3)$. Conditional on $\pi^0$, the payoff to Player 1 is at most $\delta \delta_b$. On the other hand, for each $\epsilon > 0$, the proposal $(1 - \delta \delta_b - \epsilon, \delta \delta_b + \epsilon, 0)$ is accepted by Claim 2.3. Since $s$ is an SPE, we must have $1 - \delta \delta_b - \epsilon \leq \delta \delta_b$ for each $\epsilon > 0$, and hence $1 - \delta \delta_b \leq \delta \delta_b$. Since $\pi^0 = \delta \delta_b$ by Theorem 3.2, we find $1 - \delta \delta_b \leq \delta \delta_b$, so

$$1 - 2 \delta \delta_b = 2 \sqrt{\frac{3 - \delta}{9 - 6 \delta}}.$$ Rearranging terms, we find that $\delta \geq 6 - 3\sqrt{3}$.

To prove the if part of Theorem 4.1, we construct an SPE $\hat{s}$ such that on the equilibrium path of play the proposer $\pi^0(1)$ in period 0 demands the entire surplus, his proposal is rejected, and the proposal in period 1 by Player $\pi^0(1)$ is accepted leading to an expected payoff of $\delta \delta_b$ to Player $\pi^0(1)$. Any deviation by Player $\pi^0(1)$ results in a payoff of at most $1 - \delta \delta_b$. Under the assumption that $1 - \delta \delta_b \leq \delta \delta_b$ such a deviation is not profitable. Let $\pi^0$ be the SPE provided by Theorem 3.1 with $a_1 = (1 - \delta \delta_b, \delta \delta_b, 0)$, $a_2 = (\delta \delta_b, 1 - \delta \delta_b, 0)$, and $a_3 = (\delta \delta_b, 0, 1 - \delta \delta_b)$. Strategy profile $\pi^0$ gives an expected payoff of $\delta \delta_b$ to Player 1. Similarly, let $\pi^2$ be the SPE provided by Theorem 3.1 with $a_1 = (1 - \delta \delta_b, \delta \delta_b, 0)$, $a_2 = (0, 1 - \delta \delta_b, \delta \delta_b)$, and $a_3 = (\delta \delta_b, 1 - \delta \delta_b, 0)$, and let $\pi^2$ be the SPE with $a_1 = (1 - \delta \delta_b, 0, \delta \delta_b)$, $a_2 = (0, 1 - \delta \delta_b, \delta \delta_b)$, and $a_3 = (\delta \delta_b, 0, 1 - \delta \delta_b)$. Strategy profile $\pi^2$ gives an expected payoff of $\delta \delta_b$ to Player 2 and $\pi^2$ gives an expected payoff of $\delta \delta_b$ to Player 3.

Let $\gamma$ be the SPE delivered by Theorem 3.1 with $a_1 = (1 - \delta \delta_b, \delta \delta_b, 0)$, $a_2 = (0, 1 - \delta \delta_b, \delta \delta_b)$, and $a_3 = (\delta \delta_b, 0, 1 - \delta \delta_b)$. The payoffs corresponding to $\gamma$ are $(1/3, 1/3, 1/3)$.

The one-period delay strategy profile $\hat{s}$ is defined by means of the strategy profiles $\pi^1$, $\pi^2$, and $\gamma$. The initial proposer $\pi^0(1)$ proposes to keep the entire surplus, so proposes the unit vector $e_{\pi^0(1)}$, a proposal which is rejected. For histories where the initial proposer has deviated, $\hat{s}$ is defined to be equal to $\gamma$. Otherwise, starting in period 1, $\hat{s}$ is defined to be equal to $\pi^0(1)$, where the period 0 history is deleted before applying $\pi(1)$. Whomever is
the subgame starting at history $\pi_0$.

Claim 4.4. If the first part of Theorem 4.1 follows immediately from the following

$\delta$ deviation to $x^0$. More precisely, we have the following definitions.

When the proposal $e_{\pi_1}(1)$ is made in period 0, we define $\tilde{s}$ for histories in period $t \geq 1$ by

When the proposal $x^0$ made in period 0 is not equal to $e_{\pi_1}(1)$, we define $\tilde{s}$ for histories in period $t \geq 1$ by

Let $\ell \in \{0, \ldots, t\}$. For a history $h$ in period $t$ we define "the tail" $h^{-\ell}$ of $h$ obtained by deleting the first $\ell$ periods from $h$. More precisely, for $h = (\pi_0, x^0, \ldots, x^t)$ we define $h^{-\ell} = (\pi_0, x^0, \ldots, x^\ell)$, and for $h = (\pi_0, x^0, \ldots, x^t)$ we define $h^{-\ell} = (\pi_0, x^0, \ldots, x^{t-\ell})$. In particular $h^{-0} = h$.

The claim below makes an almost trivial observation: Assume that the strategy profile $s$ is an SPE of $G$. Assume furthermore that the strategy profile $\tilde{s}$ is such that, as soon as some given history $h$ in period $t$ has been reached, the strategy profile $\tilde{s}$ is equal to $s$, where $s$ is applied after deleting the first $t$ periods from the history. Then $\tilde{s}$ is subgame perfect in the subgame starting at $h$. For a history $h \in H$, we denote the player who acts at $h$ by $\vartheta(h)$. For history $h, h'$, we write $h \succ h'$ if $h = h'$ or the history $h$ extends the history $h'$.

Claim 4.3. Let $s$ be an SPE of $G$ and let $h'$ be a history in period $t$. Let the strategy profile $\tilde{s}$ be such that $\tilde{s}_{\vartheta(h)}(h) = s_{\vartheta(h)}(h^{-\ell})$ for each history $h \succ h'$. Then $\tilde{s}$ is subgame perfect in the subgame starting at $h'$.

Proof. The subgame starting at history $h$ is isomorphic to the one starting at $h^{-\ell}$. Hence, if there would be a profitable one-shot deviation from $s$ at $h$, there would be one at $s$ at $h^{-\ell}$. □

As it is clear that $\tilde{s}$ is a strategy profile with one-period delay, the if part of Theorem 4.1 follows immediately from the following claim.

Claim 4.4. If $\delta \geq 6 - 3\sqrt{3}$, then $\tilde{s}$ is an SPE.

Proof. Consider the proposer history $(\pi^0)$. Under the strategy profile $\tilde{s}$, Player $\pi^0(1)$ makes the proposal $e_{\pi^0}(1)$, a proposal that is followed by a rejection, and the acceptance of the proposal $\tilde{s}_{\vartheta}(\pi^0, e_{\pi^0}(1), \pi_1)$ next period, leading to an expected payoff of $\delta b$ to Player $\pi^0(1)$. We argue that Player $\pi^0(1)$ does not have a profitable one-shot deviation from $\tilde{s}_{\vartheta}(\pi^0)$ at $(\pi^0)$. Consider a one-shot deviation to $x^0 \neq e_{\pi^0}(1)$, by Player $\pi^0(1)$. For histories $h \geq (\pi^0, x^0)$, it holds that $\tilde{s}_{\vartheta}(h) = \tilde{s}(h)$.

At history $(\pi^0)$, the strategy profile $\tilde{s}$ requires that the proposal $\gamma_{\pi^0}(1)$ is made and accepted, which gives Player $\pi^0(1)$ a utility equal to $1 - \delta b$. Since $\gamma$ is subgame perfect, Player $\pi^0(1)$ does not have a profitable one-shot deviation from $\gamma$, so under $\gamma$, the expected payoff of Player $\pi^0(1)$ following any proposal $y$ is less than or equal to $1 - \delta b$. Since for histories $h \geq (\pi^0, x^0)$, it holds that $\tilde{s}_{\vartheta}(h) = \gamma_{\pi^0}(h)$, we find that under $\tilde{s}$, a one-shot deviation to $x^0 \neq e_{\pi^0}(1)$ by Player $\pi^0(1)$ gives him an expected payoff less than or equal to $1 - \delta b$. Since $\delta \geq 6 - 3\sqrt{3}$, we have that $1 - \delta b \leq \delta$, so the one-shot deviation to $x^0$ is not profitable.

For all other histories, the one-shot deviation property follows from Claim 4.3. Indeed, consider $x^0 \neq e_{\pi^0(1)}$. For histories $h \geq (\pi^0, x^0)$, the one-shot deviation property follows from Claim 4.3 with $h' = (\pi^0, x^0), \ell = 0$, and $s = \gamma$. For histories $h \geq (\pi^0, e_{\pi^0(1)}, \pi^1)$, the one-shot deviation property follows from Claim 4.3 with $h' = (\pi^0, e_{\pi^0(1)}, \pi^1), \ell = 1$, and $s = \gamma_{\pi^0}(1)$. □

5. Arbitrarily long delay

In this section we show that any finite delay is compatible with an SPE, provided that the players are patient enough. By analogy with the case of one-period delay, we say that the strategy profile $s$ has $m$-period delay if, irrespective of the moves by nature, the proposals in periods 0, ..., $m - 1$ are rejected, and the proposal in period $m$ is accepted.

Let some $m \geq 1$ be given. In this section, we specify $\tilde{s} < 1$ and construct a strategy profile $\tilde{s}$ with $m$-period delay such that $\tilde{s}$ is an SPE whenever $\delta \geq \delta$. We let $\delta$ be the unique value of $\delta$ that solves the equation

$$\frac{m}{3} = 1 - \delta$$

Notice that the right-hand side of this equation is exactly $1 - \delta$. This is a decreasing function of $\delta$ with value 1 at $\delta = 0$ and value 0 at $\delta = 1$. Hence the solution exists and is unique. Moreover, for each $\delta > \delta$, it holds that $\delta^m / 3 > 1 - \delta$.

Since

$$0 = \delta^{m+2} + 2\delta^{m+1} - 3\delta^m - 27\delta + 27 \leq \delta^m + 6\delta^m - 9\delta^m - 27\delta^m + 27 = -29\delta^m + 27$$

it follows that $\delta \leq (27/29)^{1/m}$.

For $\delta \geq \delta$, the vector $(\delta^m / 3, \delta^m / 3, \delta^m / 3)$ belongs to the sets $V_1, V_2,$ and $V_3$, since $\delta^m / 3 \geq 1 - \delta b$ and $1 - \delta b > \delta b$. We let $\gamma$ be the strategy profile provided by Theorem 3.1 with $a_1 = a_2 = a_3 = (\delta^m / 3, \delta^m / 3, \delta^m / 3)$. The vector $(1/3, 1/3, 1/3)$ also belongs to the sets $V_1, V_2,$ and $V_3$. Let $\tau$ be the strategy profile provided by Theorem 3.1 with $a_1 = a_2 = a_3 = (1/3, 1/3, 1/3)$.

We construct a strategy profile $\tilde{s}$ such that on the equilibrium path of play in periods 0, ..., $m - 1$ the proposers demand the entire surplus and their proposals are rejected. If period $m$ is reached without deviations, the history of the preceding periods is deleted and as of period $m$ the strategy profile $\tau$ is followed. In particular, on the equilibrium path of play, the equal split is proposed and accepted in period $m$. However, as soon as a proposer deviates in some period $\ell < m$, the history of the preceding periods is deleted, and the strategy profile $\gamma$ is followed.

We partition the set $H$ of histories into pairwise disjoint sets $D_{-1}, D_0, \ldots, D_m$ as follows:

$$D_{-1} = \{ h \in H \mid (\pi^0, e_{\pi^0(1)}, \ldots, \pi^{m-1}, e_{\pi^{m-1}(1)} \geq h) \}$$

for $0 \leq \ell < m$.

$$D_\ell = \{ h \in H \mid \exists x^\ell \in V \setminus \{ e_{\pi^0(1)} \} \text{ such that } h \geq (\pi^0, e_{\pi^0(1)}, \ldots, \pi^{\ell-1}, e_{\pi^{\ell-1}(1)}, \pi^\ell, x^\ell) \}$$

and

$$D_m = \{ h \in H \mid h \geq (\pi^0, e_{\pi^0(1)}, \ldots, \pi^{m-1}, e_{\pi^{m-1}(1)}, \pi^m) \}.$$
The set $D_{-1}$ contains histories on the equilibrium path before period $m$. The set $D_1$ consists of histories with a deviation occurring in period $\ell < m$. Histories in the set $D_m$ involve at least $m$ periods and are without deviations in the first $m$ periods.

We define the strategy profile $s$ as follows:

\begin{equation}
\tilde{s}_{\ell}(\pi^{0},e_{\pi^{0}}(1),\ldots,\pi^{m-1},e_{\pi^{m-1}}(m)) = s_{\pi^{m}}(m) = \frac{1}{m+1}, \quad 0 \leq \ell < m.
\end{equation}

For $s\in\mathbb{R}$, we define $\tilde{s}_{\ell}(\pi^{0},e_{\pi^{0}}(1),\ldots,\pi^{m-1},e_{\pi^{m-1}}(m)) = s_{\pi^{m}}(m)$. This definition holds for $0 \leq \ell < m$.

\begin{equation}
\tilde{s}_{\ell}(\pi^{0},e_{\pi^{0}}(1)) = \tilde{s}_{\ell}(\pi^{0},e_{\pi^{0}}(1),\pi^{1}) = \tilde{s}_{\ell}(\pi^{0},e_{\pi^{0}}(1),\pi^{1},e_{\pi^{1}}(1)) = \tilde{s}_{\ell}(\pi^{0},e_{\pi^{0}}(1),\pi^{1},e_{\pi^{1}}(1),\ldots,\pi^{m-1},e_{\pi^{m-1}}(m)) = s_{\pi^{m}}(m) = \frac{1}{m+1}.
\end{equation}

We define $s_{\pi^{m}}(m)$ as the expected payoff of Player $\pi^{m}$ in period $m$.

\begin{equation}
s_{\pi^{m}}(m) = \sum_{h \in D_m} \pi^{h} \cdot \tilde{s}_{\ell}(h) = \sum_{h \in D_m} \pi^{h} \cdot \frac{1}{m+1} = \frac{1}{m+1}.\sum_{h \in D_m} \pi^{h} = \frac{1}{m+1} \cdot 1 = \frac{1}{m+1}.
\end{equation}

Eqs. (5.1)–(5.3) specify $s$ on the set of histories $D_{-1}$. Eq. (5.4) on $D_m$, and Eq. (5.5) on $D_1$ for $\ell \in \{0, \ldots, m-1\}$.

From Eq. (5.4) we get

\begin{equation}
\tilde{s}_{\ell}(\pi^{0},e_{\pi^{0}}(1),\ldots,\pi^{m-1},e_{\pi^{m-1}}(m)) = \tilde{s}_{\ell}(\pi^{0},e_{\pi^{0}}(1),\ldots,\pi^{m-1},e_{\pi^{m-1}}(m),\pi^{m}) = \frac{1}{m+1}, \quad 0 \leq \ell < m.
\end{equation}

\begin{equation}
\tilde{s}_{\ell}(\pi^{0},e_{\pi^{0}}(1),\ldots,\pi^{m-1},e_{\pi^{m-1}}(m),\pi^{m}) = \frac{1}{m+1}, \quad 0 \leq \ell < m.
\end{equation}

Theorem 5.1. Let some $m \geq 1$ be given. For every $\delta \geq \tilde{s}$, it holds that $s$ is an SPE with $m$-period delay.

\begin{proof}
We show that there are no profitable one-shot deviations from $s$.

For some $\ell \in \{0, \ldots, m-1\}$, consider the proposer history $(\pi^{0},e_{\pi^{0}}(1),\ldots,\pi^{\ell-1},e_{\pi^{\ell-1}}(\ell),\pi^{\ell})$. Under the strategy profile $s$, Player $\pi^{\ell}$ receives a payoff of $\delta/m$.

- Consider a one-shot deviation to $x^{\ell} = e_{\pi^{\ell}}(1)$. Player $\pi^{\ell}$ does not have a profitable one-shot deviation from $x^{\ell}$, $\delta/m$, and $\pi^{\ell}$ is unique in period $\ell$.

- For $\ell \in \{0, \ldots, m-1\}$, Player $\pi^{\ell}$ does not have a profitable one-shot deviation from $x^{\ell}$, $\delta/m$, and $\pi^{\ell}$ is unique in period $\ell$.

- For $\ell = 0$, consider the proposer history $(\pi^{0},e_{\pi^{0}}(1),\ldots,\pi^{m-1},e_{\pi^{m-1}}(m))$. According to $s$, both responders reject the proposer $\pi^{m}$.

- Accepting the proposal $\pi^{m}$ gives a responder $\pi^{m}$ a payoff equal to 0 and is clearly not a profitable deviation from $s$.

- For histories in $D_0, \ldots, D_m$, the one-shot deviation property follows by an application of Claim 4.3.

\end{proof}

6. Conclusion

In this paper we consider the Baron and Ferejohn (1989) model, where three players use majority voting to divide a surplus in the presence of discounting. We give a complete characterization of the set of subgame perfect equilibrium payoffs for any value of the discount factor. Our results reveal that the behavior of subgame perfect equilibrium in the BF model with three player is different from that in the case of five or more players. Indeed, in the three player case, for any given value of the discount factor there exist feasible payoff vectors that cannot arise as subgame perfect equilibrium payoffs. Nevertheless, we show that, given a vector in the interior of the set of feasible payoffs, there exists a discount factor such that the given vector is a subgame perfect equilibrium outcome.

Many papers in the experimental literature use the subgame perfect equilibrium in stationary strategies as a benchmark and conclude that this equilibrium is not supported by the data. The experimental findings are however not in contradiction to the predictions of subgame perfect equilibrium once the stationarity assumption is dropped and the discount factor is sufficiently high. Experimental evidence for a discount factor as low as 1/2 cannot be explained by non-stationary subgame perfect equilibria. Experimental papers also find significant amounts of delay before an acceptance takes place, where the delay probability increases with the discount factor. We provide necessary and sufficient conditions on the discount factor to have one period of delay, and we construct subgame perfect equilibria where arbitrary long delay occurs whenever the discount factor is sufficiently high.

Appendix. The proof of Theorem 3.1

Theorem 3.1 is proven separately in two subsections for the case $\delta \leq 3/5$ and for the case $\delta > 3/5$. In both cases the strategy profile $s$ is constructed along the same lines.

We introduce some additional notation. Consider a responder history $(\pi^{0},x^{0},\ldots,\pi^{t},x^{t})$. For some $\ell$ the strategy profile $s$ is defined recursively as follows. Set $\rho^{0}(0) = \rho^{0}$ and for each $\ell \geq 0$ let $\rho^{\ell+1}(p^{0},p^{1},x^{0},x^{1},k^{0},k^{1})$ be defined recursively as follows.

\begin{equation}
\rho^{\ell+1}(p^{0},p^{1},x^{0},x^{1},k^{0},k^{1}) = f(p^{0},x^{0},k^{0},p^{0},x^{0},k^{1}),
\end{equation}

where each of the subscripts below uses a particular specification for the function $f$. With a minor abuse of notation we write $\rho^{\ell}(h)$ rather than $\rho^{\ell}(h)$. For some $\ell \in \{0, \ldots, m-1\}$, consider the responder history $h = (\pi^{0},e_{\pi^{0}}(1),\ldots,\pi^{\ell},e_{\pi^{\ell}}(1))$. According to $s$, both responders reject the proposal $e_{\pi^{\ell}}(1)$.
The proof of Theorem 3.1 for $\delta \leq 3/5$

Throughout this subsection it is assumed that $\delta \leq 3/5$. The function $f$ is defined in Tables 1 and 2. If after the proposer history $h$ in period $t$ the proposer complies with his strategy and makes a proposal $x'_t = \rho^t(h)$, then the proposals in period $t + 1$ are determined by Table 1. If the proposer deviates from his strategy and makes a proposal $x'_t \neq \rho^t(h)$, then the proposals in period $t + 1$ are determined by Table 2. The first three columns of the tables list the six possible values for $(p^t, f^t, k^t)$, and the last three columns specify $f(p^t, x^t, f^t, k^t, p^{t+1}, y^t)$ for each of the three possible values of $p^{t+1}$.

We define the strategy profile $s$ as follows:

1. For each $t \geq 0$, for each history $h = (x^0, x^1, \ldots, x^t)$, Player $p^t$ proposes $\rho^t(h)$.
2. After history $h = (x^0, x^1, \ldots, x^t, x')$, Player $f^t$ accepts $x'$ if
   \begin{itemize}
   \item $x' = \rho^t(h)$ or
   \item $x'_j \geq \delta b$
   \end{itemize}
   and rejects otherwise.
3. After history $h = (x^0, x^1, \ldots, x^t, x')$, Player $k^t$ accepts $x'$ if
   \begin{itemize}
   \item $x' = \rho^t(h)$ and $x'_j \geq \delta b/3 + 3 \delta b$
   \item $x' \neq \rho^t(h)$ and $x'_j \geq \delta b$,
   \end{itemize}
   and rejects otherwise.

Fig. 2 shows some key thresholds as a function of $\delta$.

The intuition behind strategy profile $s$ is as follows. Table 1 determines $\rho^{t+1}(h)$ after a history $h = (x^0, x^1, \ldots, x^t, x^{t+1})$ such that $x^d = \rho^d(h)$ and $x_{d'} \geq \delta b/3 + \delta b$. In other words, it describes the proposal to be made after the equilibrium proposal was made and rejected in the previous period. According to Table 1, in period $t + 1$ Player $p^t$ obtains the best possible expected outcome, $\frac{1}{2} (1 - \delta b) + \frac{3}{2} \delta b = \frac{3}{2} \delta b$, as an encouragement for making the proposal described by $s$ in the previous period; Player $f^t$ obtains the worst possible expected outcome, $\frac{1}{2} (1 - \delta b) = \frac{1}{2} \delta b$, as a punishment for rejecting the proposal described by the strategy profile; Player $k^t$ receives the remainder, which is equal to $\frac{3}{2} \delta b$. Table 2 determines $\rho^{t+1}(h)$ after a history $h = (x^0, x^1, \ldots, x^t, x^{t+1})$ such that $x' \neq \rho^t(h)$ and $x'_{d'} \geq \delta b/3 + 3 \delta b$. In other words, it describes the proposal to be made after a deviation by the proposer in the previous period. According to Table 2, Player $p^t$ obtains an expected outcome of $\delta b + \delta b/3$, as a punishment for deviating in the previous period. Player $f^t$ obtains the best possible expected outcome of $\frac{3}{2} \delta b$, as an encouragement for rejecting a proposal that was not described by the strategy profile. Player $k^t$ receives the remainder, which is equal to $\delta b$ in expected terms.

The proof that the strategy profile $s$ defined above is an SPE proceeds by verifying that no player has a profitable one-shot deviation from $s$. The fact that $s$ is an SPE then follows from the well-known one-shot deviation principle, see Fudenberg and Tirole (1991).

The proof is divided into five steps. Claim A.1 states that $s$ has the "no delay property". Claim A.2 shows that after any responder history $h$ the responder who receives the lower share has no profitable one-shot deviation from $s$ at $h$. Claim A.3 show that after any responder history $h$ the responder who receives the higher share has no profitable one-shot deviation from $s$ at $h$. Claim A.4 show that proposals $x' \neq \rho^t(h)$ which give to the proposer more than $1 - \delta b$ are rejected by both responders. And finally, Claim A.5 shows that after any responder history $h$ the proposer has no profitable one-shot deviation from $s$ at $h$.

Claim A.1. The strategy profile $s$ has the following no delay property: if after a history $h = (x^0, x^1, \ldots, x^t)$ the players pay according to the strategy profile $s$, then $\rho^t(h)$ is proposed and accepted by one of the responders and the game ends.

Proof. The claim is immediate from the definition of $s$. Indeed, the proposer proposes $x' = \rho^t(h)$, which is accepted by Player $f^t$. □

We now turn to the absence of profitable deviations by the player who receives the smaller share.

Claim A.2. Consider a history $h = (x^0, x^1, \ldots, x^t, x')$. Player $k^t$ has no profitable one-shot deviation from $s$ at $h$.

Proof. According to the strategy profile $s$, Player $k^t$ accepts $x'$ if
\begin{itemize}
\item $x' = \rho^t(h)$ and $x'_{d'} \geq \delta b/3 + 3 \delta b$ or
\item $x' \neq \rho^t(h)$ and $x'_{d'} \geq \delta b$,
\end{itemize}
and rejects otherwise. We consider four cases:

1. $x' = \rho^t(h)$ and $x'_{d'} < \delta b/3 + \delta b$.
2. $x' = \rho^t(h)$ and $x'_{d'} \geq \delta b/3 + 3 \delta b$.
3. $x' \neq \rho^t(h)$ and $x'_{d'} < \delta b$.
4. $x' \neq \rho^t(h)$ and $x'_{d'} \geq \delta b$.

The behavior of Player $k^t$ is irrelevant if Player $f^t$ responds after Player $k^t$ and accepts. So in the following we only need to consider...
the case where Player $j$ moves before Player $k'$ or Player $j'$ moves after Player $k'$ and rejects.

Case 1. $x' = \rho'(h)$ and $x'_2 < \delta^2 b/3 + \delta b$.

According to the strategy profile $s$, Player $k'$ rejects the proposal, which leads to an expected payoff of $\delta^{t+1} (\delta b/3 + 1 - \delta b)/3 = \delta^{t+1} (\delta b/3 + b) = \delta^{t+1} (\delta^2 b/3 + \delta b)$. If Player $k'$ deviates and accepts, this gives him a payoff less than $\delta^{t+1} (\delta^2 b/3 + \delta b)$. So acceptance is not a profitable deviation.

Case 2. $x' = \rho'(h)$ and $x'_2 \geq \delta^2 b/3 + \delta b$.

According to the strategy profile $s$, Player $k'$ accepts the proposal, which leads to a payoff of at least $\delta^{t+1} (\delta^2 b/3 + \delta b)$. If Player $k'$ deviates and rejects, the expected payoff is $\delta^{t+1} (\delta^2 b/3 + \delta b)$. So rejection is not a profitable deviation.

Case 3. $x' \neq \rho'(h)$ and $x'_2 < \delta b$.

According to the strategy profile $s$, Player $k'$ rejects the proposal, leading to an expected payoff of $\delta^{t+1} (1 - \delta b)/3 = \delta^{t+1} b$. If Player $k'$ deviates and accepts, it leads to a payoff less than $\delta^{t+1} b$. So acceptance is not a profitable deviation.

The behavior of Player $j'$ is irrelevant if Player $k'$ responds after Player $j'$ and accepts. So in the following we only need to consider the case where Player $k'$ moves before Player $j'$ or Player $k'$ moves after Player $j'$ and rejects.

Case 1. $x' = \rho'(h)$.

Note that if $x' = \rho'(h)$, then $x'_2 \geq \delta b$. Indeed, if $x' = a$, then $x'_2 \geq \delta b$. Moreover, any proposal described by Tables 1 and 2 satisfies $x'_2 = \delta b$ or $x'_2 = \delta b$. According to the strategy profile $s$, Player $j'$ accepts, which leads to a payoff of $\delta x'_2 \geq \delta^2 t$. If Player $j'$ rejects $x'$, this leads to an expected payoff of $\delta^{t+1} (1 - \delta b)/3 = \delta^{t+1} b$. So rejection is not a profitable deviation.

Case 2. $x' \neq \rho'(h)$ and $x'_2 \geq \delta b$.

According to the strategy profile $s$, Player $j'$ accepts, which leads to a payoff of at least $\delta^{t+1} b$. If Player $j'$ deviates and rejects $x'$, this leads to an expected payoff of $\delta^{t+1} (1 - \delta b)/3 + 2 \delta^2 b/3 = \delta^{t+1} b$. So rejection is not a profitable deviation.

Case 3. $x' \neq \rho'(h)$ and $x'_2 < \delta b$.

According to the strategy profile $s$, Player $j'$ rejects, which leads to an expected payoff of $\delta^{t+1} (1 - \delta b)/3 + 2 \delta^2 b/3 = \delta^{t+1} b$. If Player $j'$ deviates and accepts, he receives a payoff less than or equal to $\delta^{t+1} b$. So acceptance is not a profitable deviation.

We have checked that responders have no profitable one-shot deviations. We now turn our attention to the proposer. The following claim states that whenever a proposer deviates from $s$ and demands a share larger than $1 - \delta b$, the proposal will be rejected.

**Claim A.4.** Consider a history $h = (\pi^0, x^0, \ldots, \pi^t)$. A proposal $x' \neq \rho'(h)$ such that $x'_{2'} > 1 - \delta b$ is rejected by both responders.

**Proof.** Note that $x'_2 + x'_2 < \delta b$. Since by definition Player $k'$ is the responder who is offered the smaller share, we have $x'_2 < \delta b/2$. Moreover, the assumption $\delta < 3/5$ implies $\delta b/2 \leq \delta b$, thus $x'_2 < \delta b$. It follows by the definition of $s$ that Player $k'$ rejects the proposal $x'$. Player $j'$ rejects the proposal $x'$ since $x'_2 < \delta b$. □

The next claim states that a proposer does not have a profitable one-shot deviation.

**Claim A.5.** Consider a history $h = (\pi^0, x^0, \ldots, \pi^t)$. Player $p'$ has no profitable one-shot deviation from $s$ at $h$.

**Proof.** If $\rho'(h) = a$, then $\sigma_{a'} \geq 1 - \delta b$. Moreover, any proposal $x$ described in both Tables 1 and 2 satisfies $x_{d'} = 1 - \delta b$ or $x_{p'} = 1 - \delta b$. Since the proposal $\rho'(h)$ is accepted, it follows that following $s$ leads to a payoff greater than or equal to $\delta^t (1 - \delta b)$ for Player $p'$. If Player $p'$ proposes $x' \neq \rho'(h)$ such that $x'_{2'} > 1 - \delta b$, then the proposal is rejected by Claim A.4 and the expected payoff for Player $p'$ is equal to $\delta^{t+1} (1 - \delta b)/3 + \delta b/3 \leq \delta^{t+1} (1 - \delta b)$. If Player $p'$ proposes $x' \neq \rho'(h)$ such that $x_{p'} \leq 1 - \delta b$, then the proposal is either accepted and leads to a payoff of $\delta x'_{p'} \leq \delta^t (1 - \delta b)$, or rejected and leads to an expected payoff of $\delta^{t+1} (1 - \delta b)/3 + \delta b/3 \leq \delta^{t+1} (1 - \delta b)$. In both cases, the payoff of Player $p'$ is less than or equal to $\delta^t (1 - \delta b)$. □

**The proof of Theorem 3.1 for $\delta > 3/5$**

Throughout this subsection it is assumed that $\delta > 3/5$. The proof of the previous subsection breaks down, since Claim A.4 is no longer true. If $\delta > 3/5$, it is no longer guaranteed that a proposal $x' \neq \rho'(h)$ such that $x'_{2'} > 1 - \delta b$ is rejected by responder $k'$, since it is not necessarily the case that $x'_2 \leq \delta b$. We therefore change the construction of the strategy profile $s$ and replace Table 2 by Table 3. As before, Table 1 gives the values for $f(p', x', x', k', p', p', y')$ when $x' = y'$ and Table 3 when $x' \neq y'$. For $x'_2 \in [0, 1/2]$, $\theta(x'_2)$ measures the excess of $x'_2$ over $\delta b$ when $\delta b < x'_2 < \delta b/2$ and is equal to zero otherwise. Notice that when $\delta > 3/5$, it holds that $\delta b < \delta b/2$. More precisely, $\theta(x'_2)$ is defined by

$$\theta(x'_2) = \begin{cases} x'_2 - \delta b, & \text{if } \delta b < x'_2 < \delta b/2, \\ 0, & \text{otherwise.} \end{cases}$$

(2.1)

If $\theta(x'_2) = 0$, then Table 3 is identical to Table 2. In case $\theta(x'_2) \neq 0$, Player $j'$ obtains an expected payoff of $\tilde{b} - \theta(x'_2)$. Player $j'$ is still rewarded for rejecting a proposal that was not described by the strategy profile, but he shares a part of the reward with Player $k'$, who receives the remainder $x'_2/\delta$. This makes Player $k'$ indifferent between accepting and rejecting $x'_2$.

We define the strategy profile $s$ as follows:

1. For each $t \geq 0$, for each history $h = (\pi^0, x^0, \ldots, \pi^t)$, Player $p'$ proposes $\rho'(h)$.
2. After history $h = (\pi^0, x^0, \ldots, \pi^t, x')$, Player $j'$ accepts $x'$ if
   - $[x' = \rho'(h)]$ or
   - $[x'_2 \leq \delta b$ and $x'_2 \geq \delta b]$ or
   - $[\delta b < x'_2 < \delta b/2$ and $x'_2 \geq \delta b + \delta b - x'_2]$ or
   - $[x'_2 \geq \delta b/2$ and $x'_2 \geq \delta b]$, and rejects otherwise.
3. After history \( h = (\pi^0, x^0, \ldots, \pi^t, x^t) \), Player \( k^t \) accepts \( x^t \) if
\[ |x^t - \rho^t(h)| < \delta b^t /3 + \delta b \] or \[ |x^t - \rho^t(h)| < \delta b^t /2 \], and rejects otherwise.

The proof of the fact that the strategy profile \( s \) described above is the SPE is divided into six steps, five of which are the same as in the proof for \( \delta \leq 3/5 \). We need one additional step to provide a bound on \( \theta(x^t_k) \).

**Claim A.6** states that \( s \) has the “no delay property”.

**Claim A.6.** The strategy profile \( s \) has the following no delay property: if after a history \( h = (\pi^0, x^0, \ldots, \pi^t) \) the players play according to the strategy profile \( s \), then \( \rho^t(h) \) is proposed and accepted by one of the responders and the game ends.

**Proof.** The claim is immediate from the definition of \( s \). Indeed, the proposer proposes \( x^t = \rho^t(h) \), which is accepted by Player \( j^t \).  \( \square \)

**Claim A.7** shows that, after any responder history \( h \), the responder with the smaller share, \( k^t \), has no profitable one-shot deviations from \( s \) at \( h \).

**Claim A.7.** Consider a history \( h = (\pi^0, x^0, \ldots, \pi^t, x^t) \). Player \( k^t \) has no profitable one-shot deviation from \( s \) at \( h \).

**Proof.** According to the strategy profile \( s \), Player \( k^t \) accepts \( x^t \) if
\[ |x^t - \rho^t(h)| < \delta b^t /3 + \delta b \] or \[ |x^t - \rho^t(h)| < \delta b^t /2 \], and rejects otherwise. We consider five cases:

1. \( x^t = \rho^t(h) \) and \( x^t \leq \delta b^t /3 + \delta b \) - the proof is the same as in **Claim A.2**.
2. \( x^t = \rho^t(h) \) and \( x^t \geq \delta b^t /3 + \delta b \) - the proof is the same as in **Claim A.2**.
3. \( x^t \neq \rho^t(h) \) and \( x^t \leq \delta b \) - the proof is the same as in **Claim A.2**.
4. \( x^t \neq \rho^t(h) \) and \( \delta b < x^t < \delta b^t /2 \)
5. \( x^t \neq \rho^t(h) \) and \( x^t \geq \delta b^t /2 \).

The behavior of Player \( k^t \) is irrelevant if Player \( j^t \) responds after Player \( k^t \) and accepts. So in the following we only need to consider the case where Player \( j^t \) moves before Player \( k^t \) or Player \( j^t \) moves after Player \( k^t \) and rejects.

Case 4. \( x^t \neq \rho^t(h) \) and \( x^t \leq \delta b^t /2 \).

According to the strategy profile \( s \), Player \( k^t \) rejects the proposal. Recalling that \( \theta(x^t_k) = x^t_k / \delta - \delta \) and using Eq. \( (3.1) \), the resulting expected payoff is equal to \( \delta^t + \frac{1}{2}(1 - \delta + \theta(x^t_k)) + \frac{1}{2}(1 - \delta + \theta(x^t_j)) = \delta^t + \frac{1}{2}(1 - \delta + \theta(x^t_k) + \theta(x^t_j)) = \delta^t + b + \theta(x^t_j) = \delta^t + b \).

If Player \( k^t \) deviates and accepts, he obtains a payoff of \( \delta^t + b \). So acceptance is not a profitable deviation.

Case 5. \( x^t \neq \rho^t(h) \) and \( x^t \geq \delta b^t /2 \).

According to the strategy profile \( s \), Player \( k^t \) accepts the proposal, which leads to a payoff of at least \( \delta b^t /2 \). If Player \( k^t \) deviates and rejects, he obtains an expected payoff of \( \delta^t + \frac{1}{2}(1 - \delta + \theta(x^t_k)) = \delta^t + \frac{1}{2}(1 - \delta + \theta(x^t_j)) = \delta^t + b + \theta(x^t_j) = \delta^t + b + \delta b /3 < \delta b^t /2 \) for \( \delta > 3/5 \). So rejection is not a profitable deviation. \( \square \)

**Claim A.8** puts a bound on \( \theta(x^t_k) \).

**Claim A.8.** For every \( x^t_k \in [0, 1/2] \), it holds that \( \theta(x^t_k) < \delta(\delta - b) \).

**Proof.** It follows from the definition that \( \theta(x^t_k) < \frac{1}{2} \delta - b \). Hence
\[ \delta b - \theta(x^t_k) > \delta b - \frac{1}{2} \delta - b = (\delta - \frac{1}{2}) \delta b + b > (\delta - \frac{1}{2}) \delta b + b \]
\[ = 2b b > \delta b \]

where we have used the fact that \( \delta > 3/5 \) and \( b > 2b > 0 \).  \( \square \)

**Claim A.9** shows that, after any responder history \( h \), the responder with the higher share, \( j^t \), has no profitable one-shot deviation from \( s \) at \( h \).

**Claim A.9.** Consider a history \( h = (\pi^0, x^0, \ldots, \pi^t, x^t) \). Player \( j^t \) has no profitable one-shot deviation from \( s \) at \( h \).

**Proof.** According to the strategy profile \( s \), Player \( j^t \) accepts \( x^t \) if
\[ |x^t - \rho^t(h)| \]
\[ |x^t \leq \delta b | \]
\[ |x^t \geq \delta b^t /2 | \]

and rejects otherwise. We consider five cases:

1. \( x^t = \rho^t(h) \).
2. \( x^t \neq \rho^t(h) \) and \( x^t \leq \delta b \) or \( x^t \geq \delta b^t /2 \) and \( x^t \geq \delta b - x^t_k \) - the proof is the same as in **Claim A.3**.
3. \( x^t \neq \rho^t(h) \) and \( x^t \leq \delta b \) or \( x^t \geq \delta b^t /2 \) and \( x^t \geq \delta b - x^t_k \) - the proof is the same as in **Claim A.3**.
4. \( x^t \neq \rho^t(h) \) and \( \delta b < x^t < \delta b^t /2 \)
5. \( x^t \neq \rho^t(h) \) and \( \delta b < x^t < \delta b^t /2 \) and \( x^t \geq \delta b - x^t_k \).

The behavior of Player \( j^t \) is irrelevant if Player \( k^t \) responds after Player \( j^t \) and accepts. So in the following we only need to consider the case where Player \( k^t \) moves before Player \( j^t \) or Player \( k^t \) moves after Player \( j^t \) and rejects.

Case 1. \( x^t = \rho^t(h) \).

First we argue that if \( x^t = \rho^t(h) \), then \( x^t \geq \delta b \). If \( x^t = a \), then \( x^t \geq \delta b \). Moreover, any proposal in **Table 3** satisfies \( x^t_k = \delta b \) or \( x^t_j = \delta b \). And any proposal in **Table 3** offers the responders the shares \( \delta b \) and \( \theta(x^t_k) \), or the shares \( \delta b - \theta(x^t_k) \) and \( \delta b - \theta(x^t_j) \). In view of **Claim A.8**, in the case where a player offers at least \( \delta b \), so \( x^t \geq \delta b \).

According to the strategy profile \( s \), Player \( j^t \) accepts, leading to a payoff \( \delta^t + \delta b \). If Player \( j^t \) deviates and rejects \( x^t \), this leads to an expected payoff of \( \delta^t + \delta b /3 = \delta^t + \delta b /2 \). So rejection is not a profitable deviation.

Case 2. \( x^t \neq \rho^t(h) \) and \( |x^t_k \leq \delta b | \)

According to the strategy profile \( s \), Player \( j^t \) rejects, leading to an expected payoff of \( \delta^t + \delta b /3 = \delta^t + \delta b /2 \). If Player \( j^t \) deviates and accepts, he receives a payoff less than \( \delta^t + \delta b /2 \). So acceptance is not a profitable deviation.
Case 4. \( x' \neq \rho'(h) \) and \( \delta b < x'_k < \delta b/2 \) and \( x'_k \geq \delta b - x'_k \).

According to the strategy profile \( s \), Player \( j' \) accepts, leading to a payoff of at least \( \delta^t + \delta^t b + \delta^t b - \delta^t x'_k \). If Player \( j' \) deviates and rejects, this leads to an expected payoff of

\[
\delta^{t+1}(1 - \delta b - \theta(x_{k+1}^{-1})) + \frac{2}{3}(\delta b - \theta(x_{k+1}^{-1})) = \delta^{t+1}(1 - \delta b - \theta(x_{k+1}^{-1})) - \delta^t x'_k,
\]

where we used (3.1) and the fact that \( \theta(x_k') = x_k' / \delta - b \). So rejection is not a profitable deviation.

Case 5. \( x' \neq \rho'(h) \) and \( \delta b < x'_k < \delta b/2 \) and \( x'_k < \delta b - \delta x'_k \).

According to the strategy profile \( s \), Player \( j' \) rejects, leading to an expected payoff of \( \delta^t + \delta^t b + \delta^t b - \delta^t x'_k \). If Player \( j' \) deviates and accepts, this gives him a payoff of at most \( \delta^t + \delta^t b + \delta^t b - \delta^t x'_k \). So acceptance is not a profitable deviation.

Claim A.10 shows that a proposal \( x' \neq \rho'(h) \) which gives to the proposer more than \( 1 - \delta b \) is rejected by both responders.

Claim A.10. Consider a history \( h = (\pi^0, x^0, \ldots, x^t) \). A proposal \( x' \neq \rho'(h) \) such that \( x'_{p^*} > 1 - \delta b \) is rejected by both responders.

**Proof.** Note that \( x'_k + x'_k < \delta b \). Since \( k' \) is the responder who is offered the smaller share, we have \( x'_k < \delta b/2 \). According to the strategy profile \( s \), Player \( k' \) rejects the proposal. Consider player \( j' \).

If \( x'_k \leq \delta b \) then, according to the strategy profile \( s \), Player \( j' \) rejects the proposal as \( x'_k < \delta b \). If \( \delta b < x'_k < \delta b/2 \) then, according to the strategy profile \( s \), Player \( j' \) rejects the proposal as \( x'_k - 1 - \delta b - x'_k = \delta b - x'_k < \delta b + \delta b - x'_k. \)

Claim A.11 shows that after any proposer history, the proposer has no profitable one-shot deviation from \( s \) at \( h \).

Claim A.11. Consider a history \( h = (\pi^0, x^0, \ldots, x^t) \). Player \( p^* \) has no profitable one-shot deviation from \( s \) at \( h \).

**Proof.** First we argue that following \( s \) leads to a payoff of at least \( \delta^t (1 - \delta b) \) for Player \( p^* \). If \( \rho'(h) = a, \) then \( a_{p^*} \geq 1 - \delta b. \) Otherwise, according to \( s, \) the proposer demands one of the four shares: \( 1 - \delta b, 1 - \delta b + \theta(x_{k+1}^{-1}), \) and \( 1 - \delta b - \theta(x_{k+1}^{-1}). \) Each of these is at least \( 1 - \delta b, \) where the fact that \( 1 - \delta b - \theta(x_{k+1}^{-1}) \) is at least \( 1 - \delta b \) follows from Claim A.8.

If Player \( p^* \) proposes \( x' \neq \rho'(h) \) such that \( x'_{p^*} > 1 - \delta b \), then the proposal is rejected according to Claim A.10 and the expected payoff for Player \( p^* \) is equal to \( \delta^t + \delta^t b + \delta^t b/3 \leq \delta^t (1 - \delta b) \).

If Player \( p^* \) proposes \( x' \neq \rho'(h) \) such that \( x'_{p^*} \leq 1 - \delta b \), then the proposal is either accepted and leads to a payoff of \( \delta x'_{p^*} \leq \delta^t (1 - \delta b)/3 + \delta b/3 \leq \delta^t (1 - \delta b) \), or is rejected and leads to an expected payoff of \( \delta^t + (1 - \delta b)/3 + \delta b/3 \leq \delta^t (1 - \delta b). \)